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Collins' and Eisner's algorithms

Syntactic analysis (5LN455)

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Based on slides by Marco Kuhlmann

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Collins' algorithm

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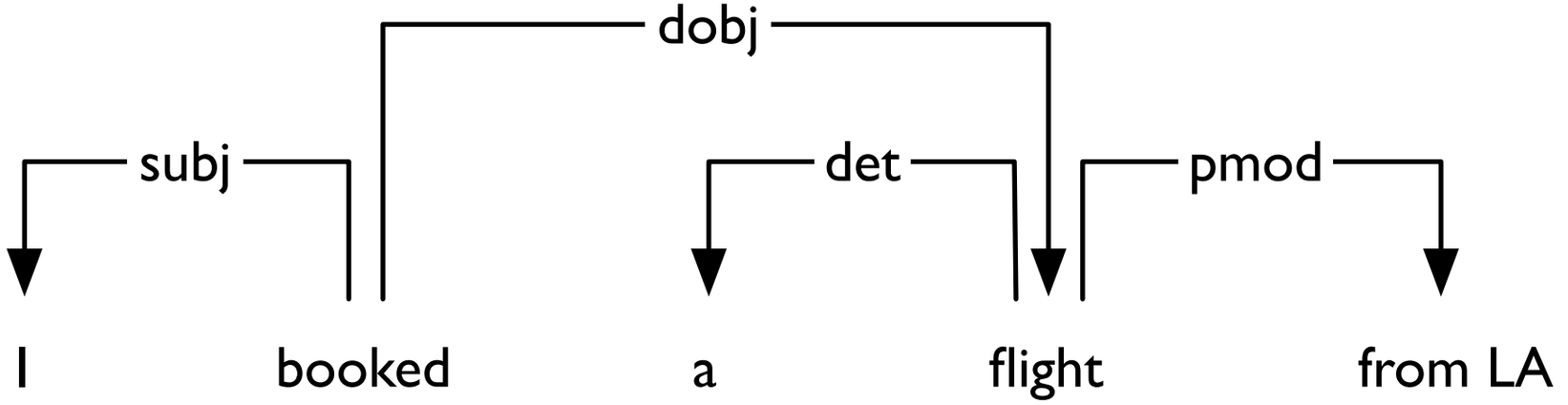
Recap: Dependency trees



- In an arc $h \rightarrow d$, the word h is called the head, and the word d is called the dependent.
- The arcs form a rooted tree.

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Recap: The arc-factored model

- To score a dependency tree, score the individual arcs, and combine the score into a simple sum.

  $\text{score}(t) = \text{score}(a_1) + \ldots + \text{score}(a_n)$

- Define the score of an arc $h \rightarrow d$ as the weighted sum of all features of that arc:

  $\text{score}(h \rightarrow d) = f_1 w_1 + \ldots + f_n w_n$

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Collin's algorithm

- Collin's algorithm is a simple algorithm for computing the highest-scoring dependency tree under an arc-factored scoring model.
- It can be understood as an extension of the CKY algorithm to dependency parsing.
- Like the CKY algorithm, it can be characterized as a bottom-up algorithm based on dynamic programming.

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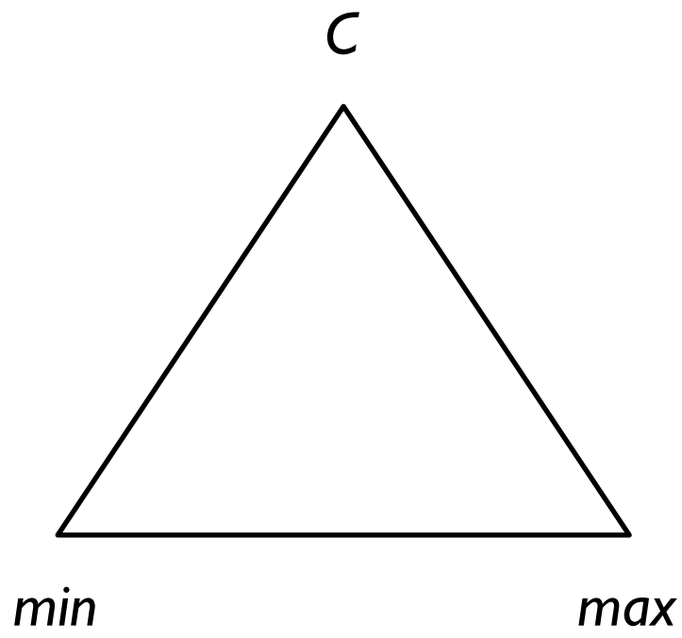
Collins' algorithm

Signatures



[*min*, *max*, *C*]

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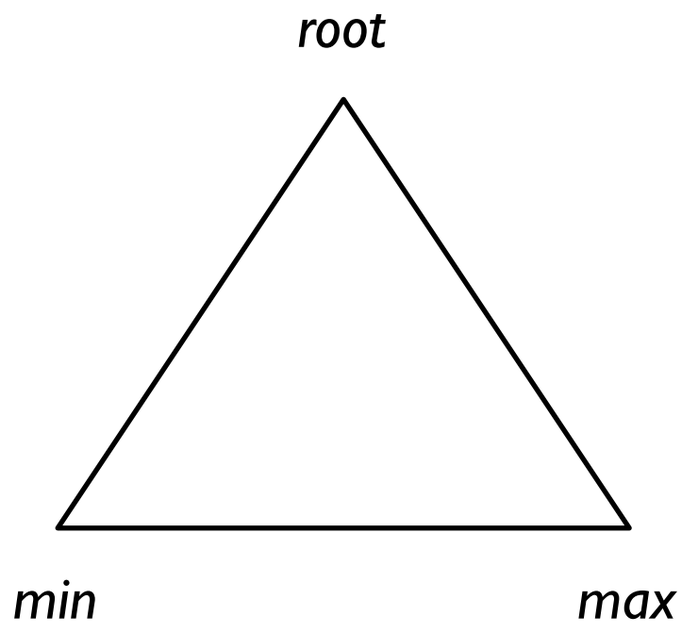
Collins' algorithm

Signatures



[*min*, *max*, *root*]

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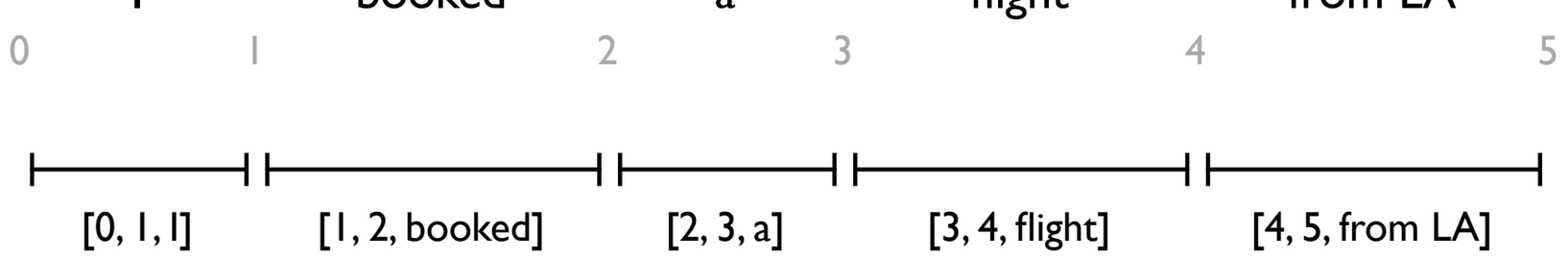
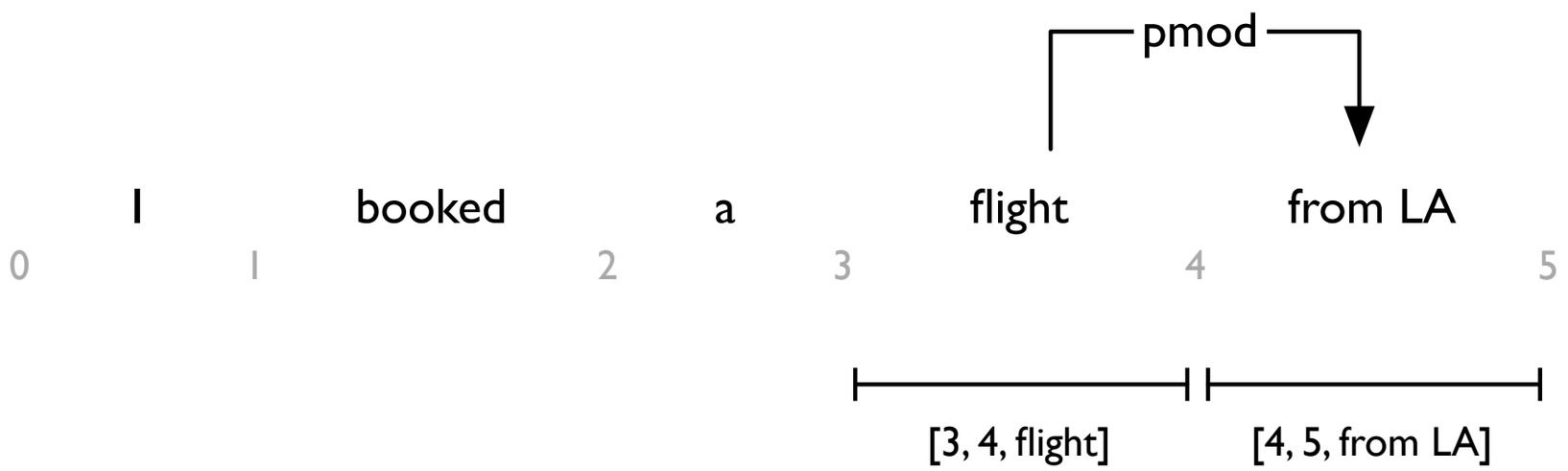
Collins' algorithm

Initialization



Adding a left-to-right arc



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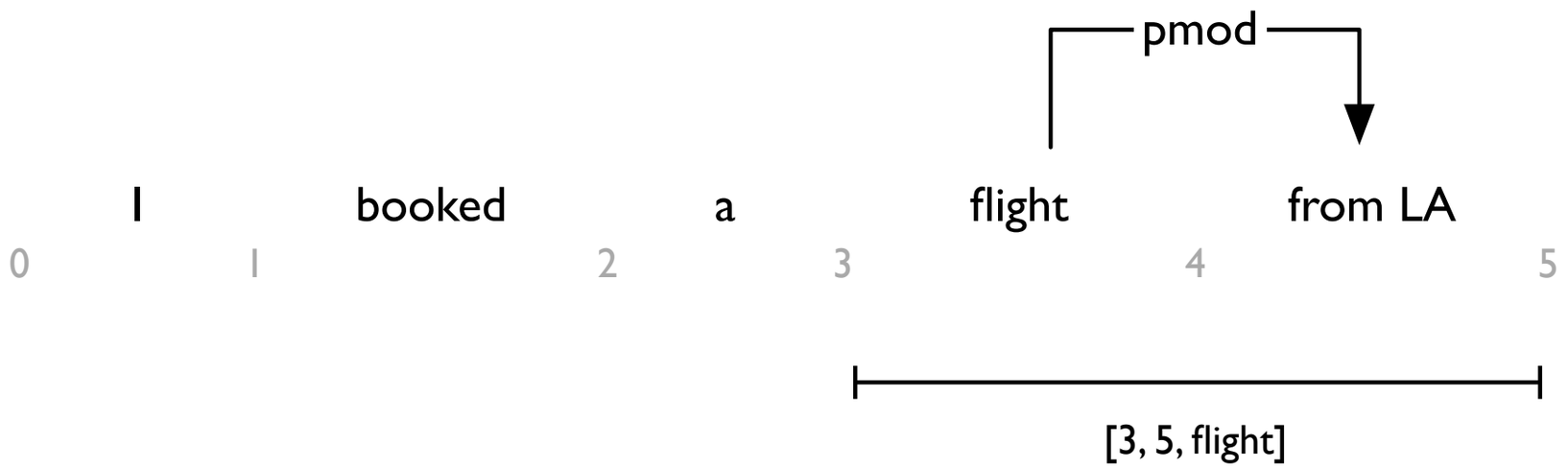

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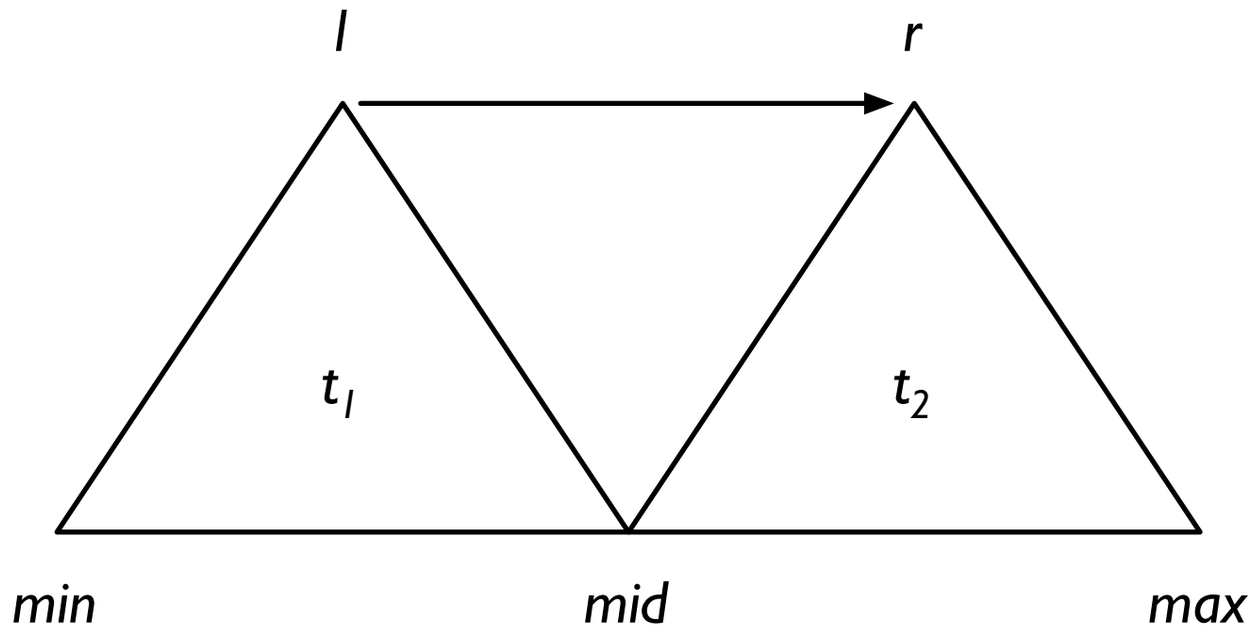

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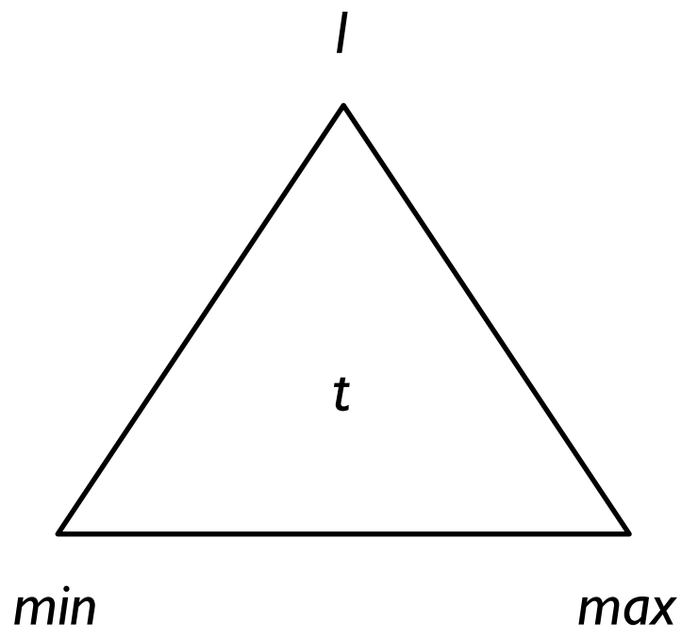
Adding a left-to-right arc

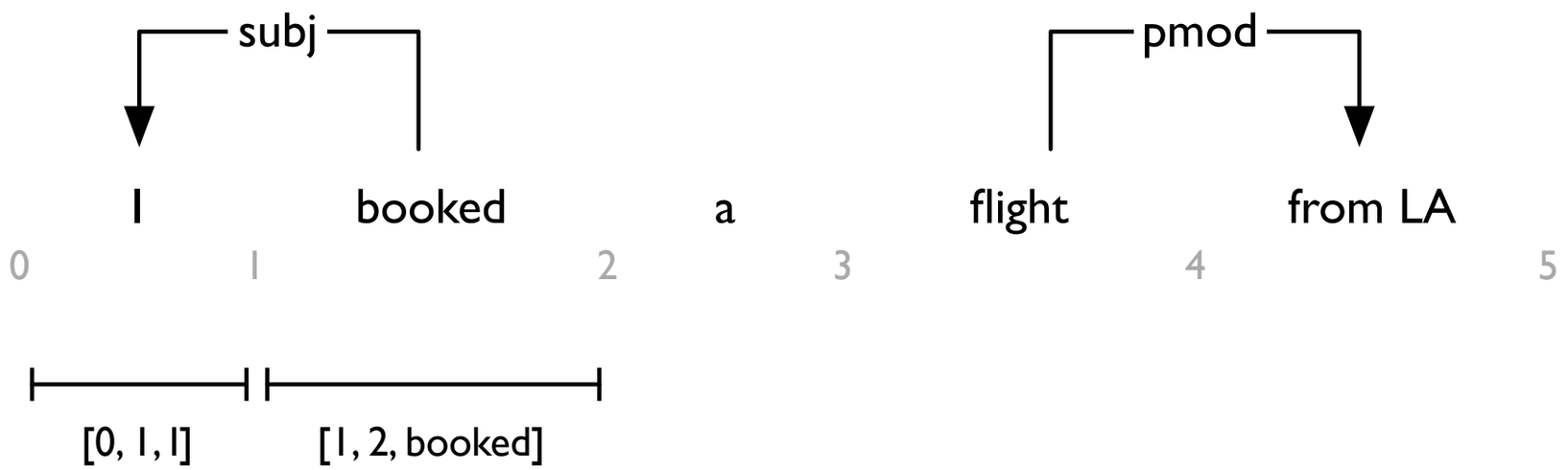


$$\text{score}(t) = \text{score}(t_1) + \text{score}(t_2) + \text{score}(l \rightarrow r)$$

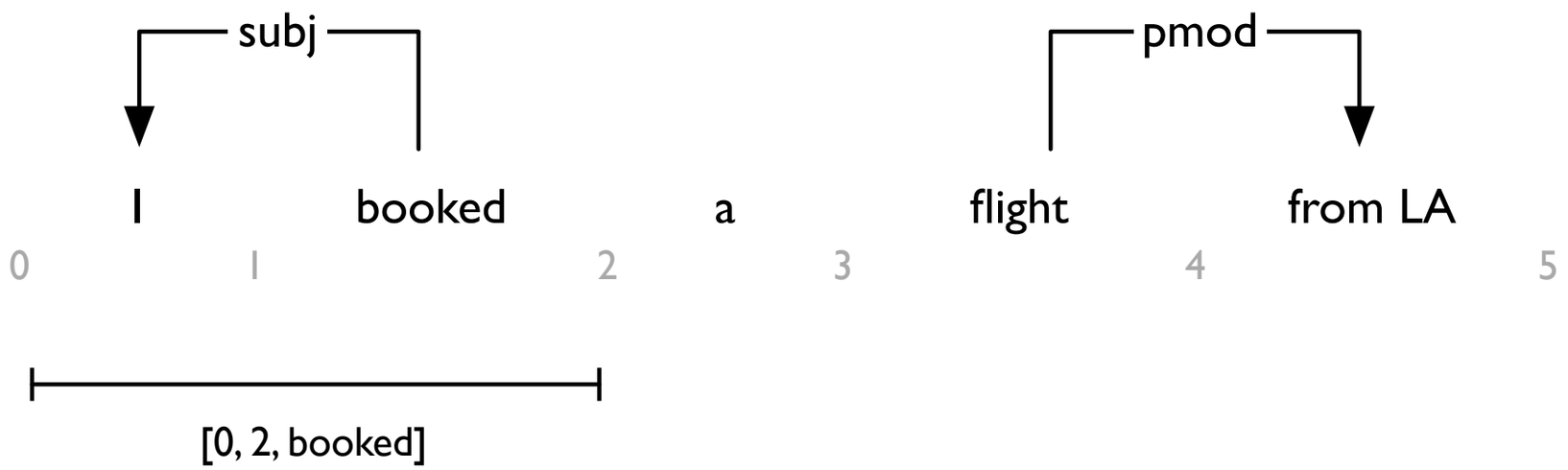
Adding a left-to-right arc

```
for each [min, max] with max - min > 1 do
  for each l from min to max - 2 do
    double best = score[min][max][l]
    for each r from l + 1 to max - 1 do
      for each mid from l + 1 to r do
        t1 = score[min][mid][l]
        t2 = score[mid][max][r]
        double current = t1 + t2 + score(l → r)
        if current > best then
          best = current
    score[min][max][l] = best
```

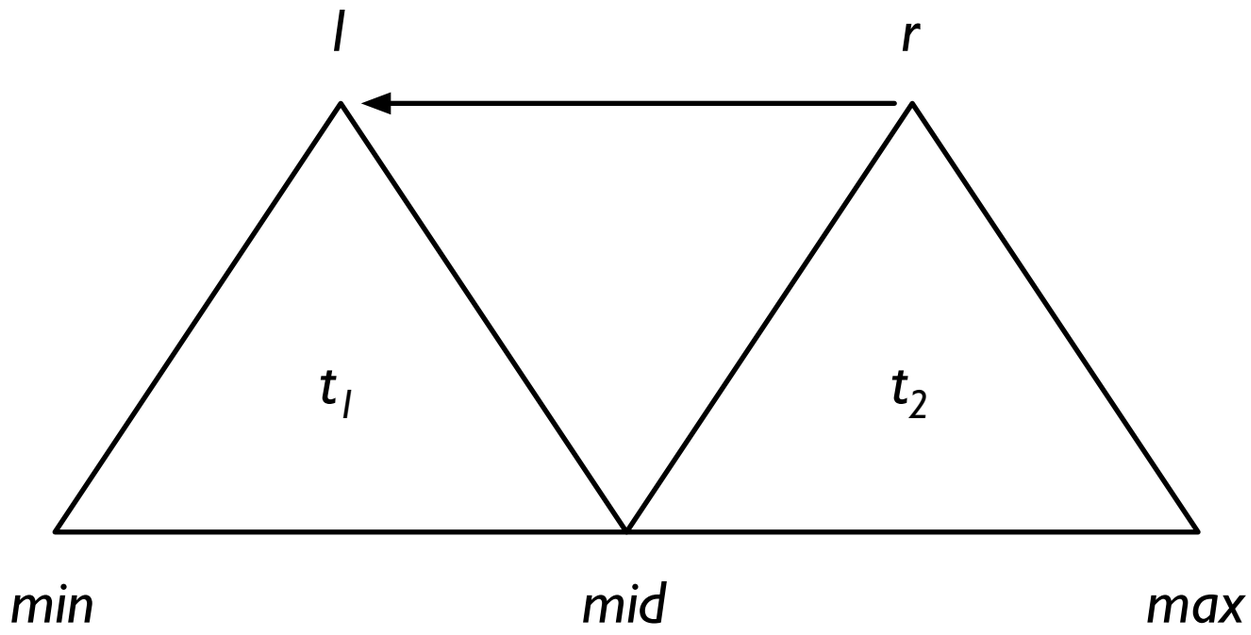
Adding a right-to-left arc



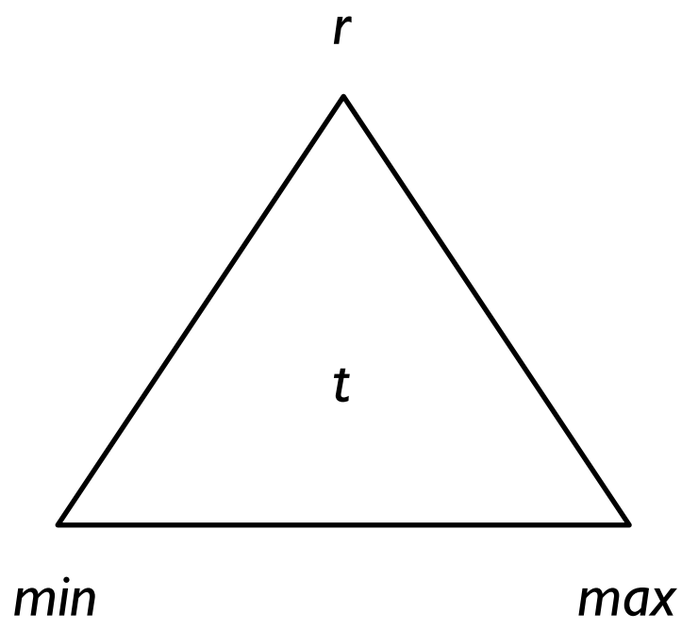
Adding a right-to-left arc



Adding a right-to-left arc



Adding a right-to-left arc

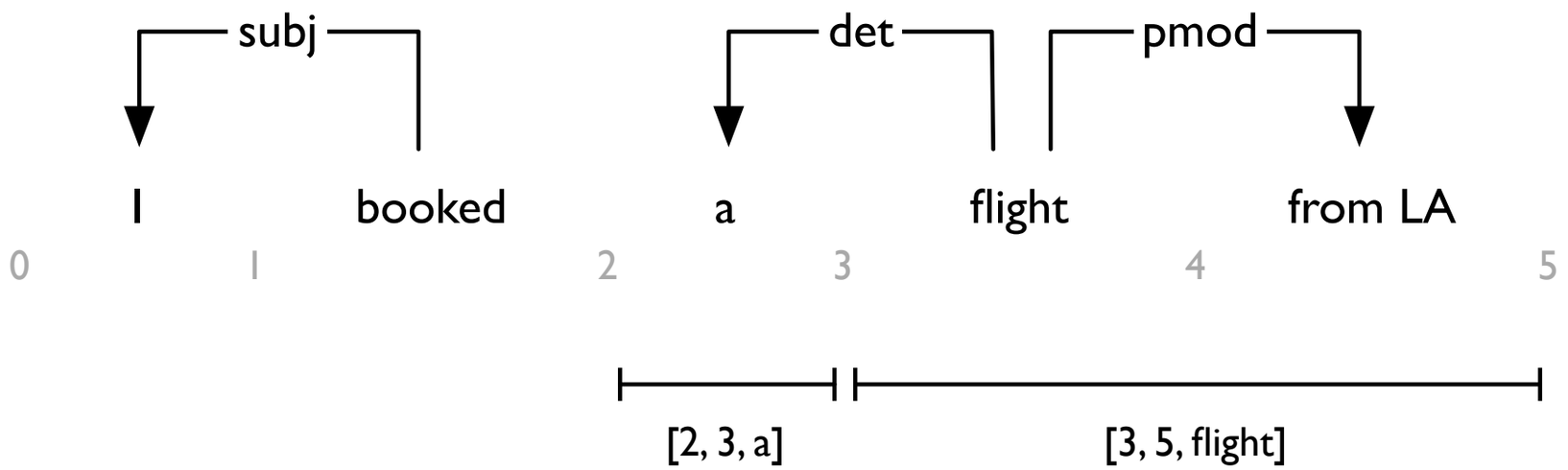


$$\text{score}(t) = \text{score}(t_1) + \text{score}(t_2) + \text{score}(r \rightarrow l)$$

Adding a right-to-left arc

```
for each [min, max] with max - min > 1 do
  for each r from min + 1 to max - 1 do
    double best = score[min][max][r]
    for each l from min to r - 1 do
      for each mid from l + 1 to r do
        t1 = score[min][mid][l]
        t2 = score[mid][max][r]
        double current = t1 + t2 + score(r → l)
        if current > best then
          best = current
    score[min][max][r] = best
```

Finishing up



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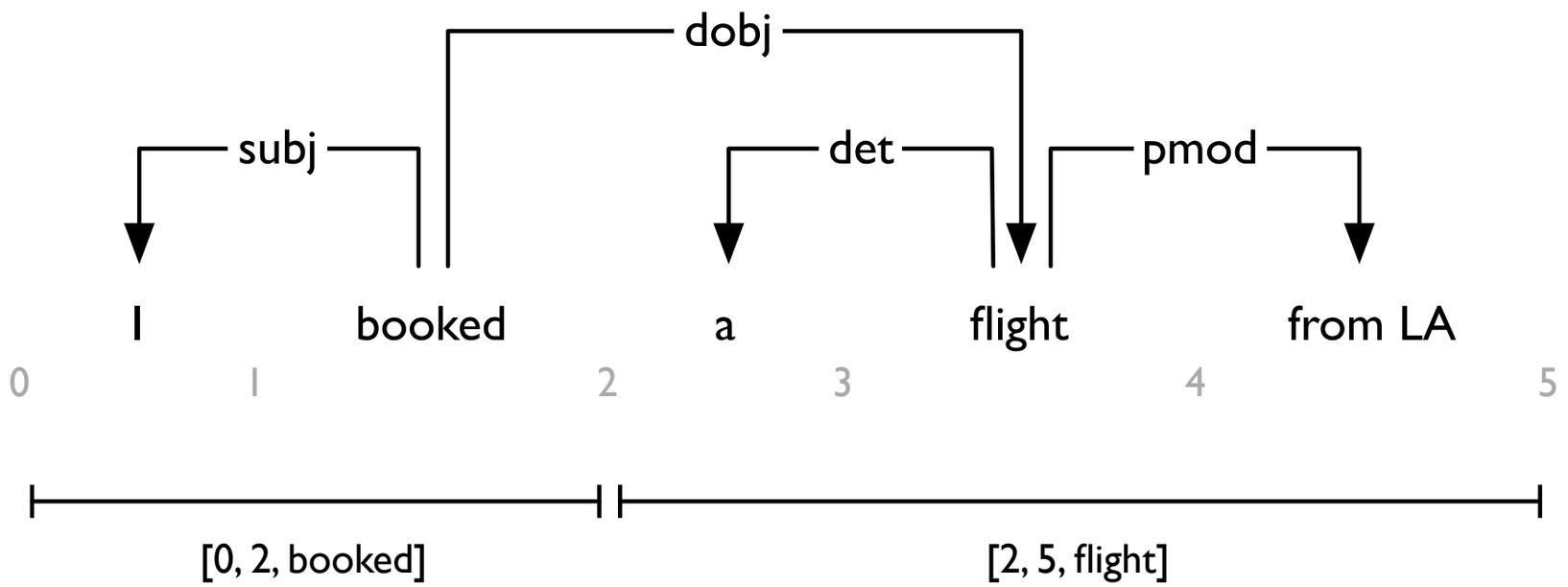

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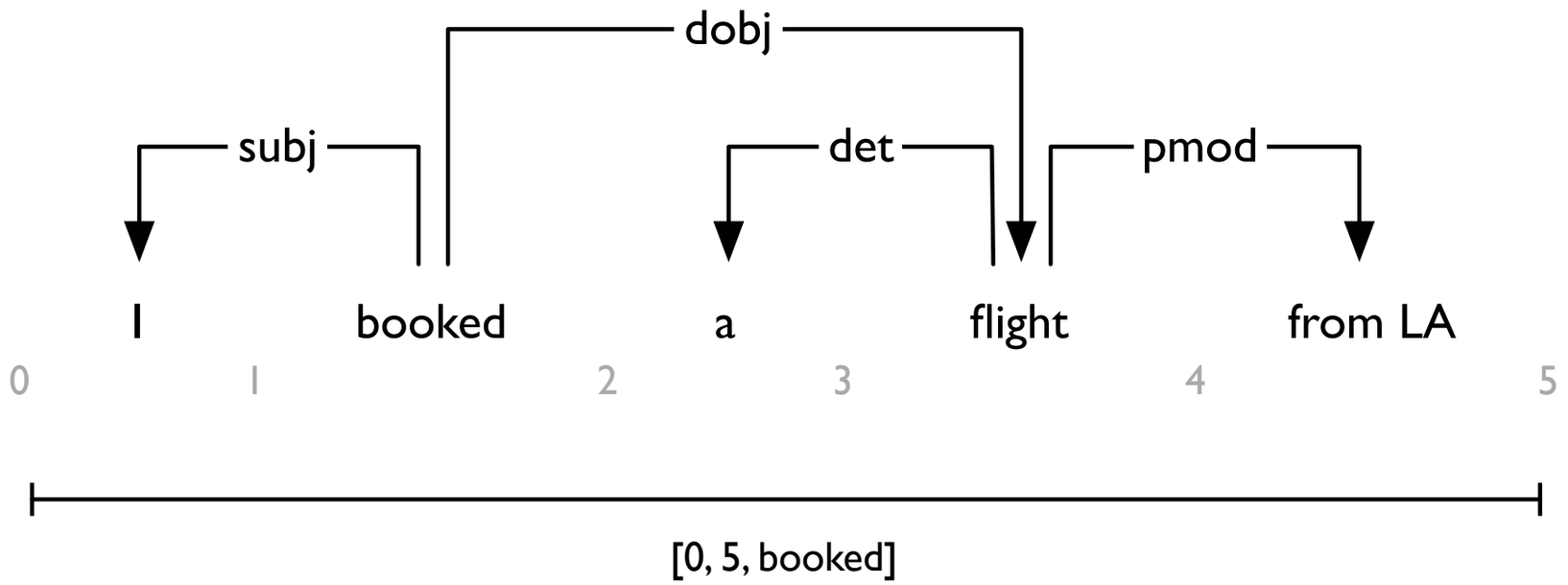

Complexity analysis

- Space requirement:
  $O(|w|^3)$
- Runtime requirement:
  $O(|w|^5)$

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Extension to the labeled case

- It is important to distinguish dependencies of different types between the same two words.

  Example: subj, dobj

- For this reason, practical systems typically deal with labeled arcs.
- The question then arises how to extend Collins' algorithm to the labeled case.

Naive approach

- Add an innermost loop that iterates over all edge labels in order to find the combination that maximizes the overall score.
- For each step of the original algorithm, we now need to make $|L|$ steps, where L is the set of all labels.

Smart approach

- Before parsing, compute a table that lists, for each head–dependent pair (h, d), the label that maximizes the score of arcs $h \rightarrow d$.
- During parsing, simply look up the best label in the precomputed table.
- This adds (not multiplies!) a factor of $|L||w|^2$ to the overall runtime of the algorithm.

Summary

- Collins' algorithm is a CKY-style algorithm for computing the highest-scoring dependency tree under an arc-factored scoring model.
- It runs in time $O(|w|^5)$.
  This may not be practical for long sentences.

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Eisner's algorithm

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Eisner's algorithm

- With its runtime of $O(|w|^5)$, Collins' algorithm may not be of much use in practice.
- With Eisner's algorithm we will be able to solve the same problem in $O(|w|^3)$.

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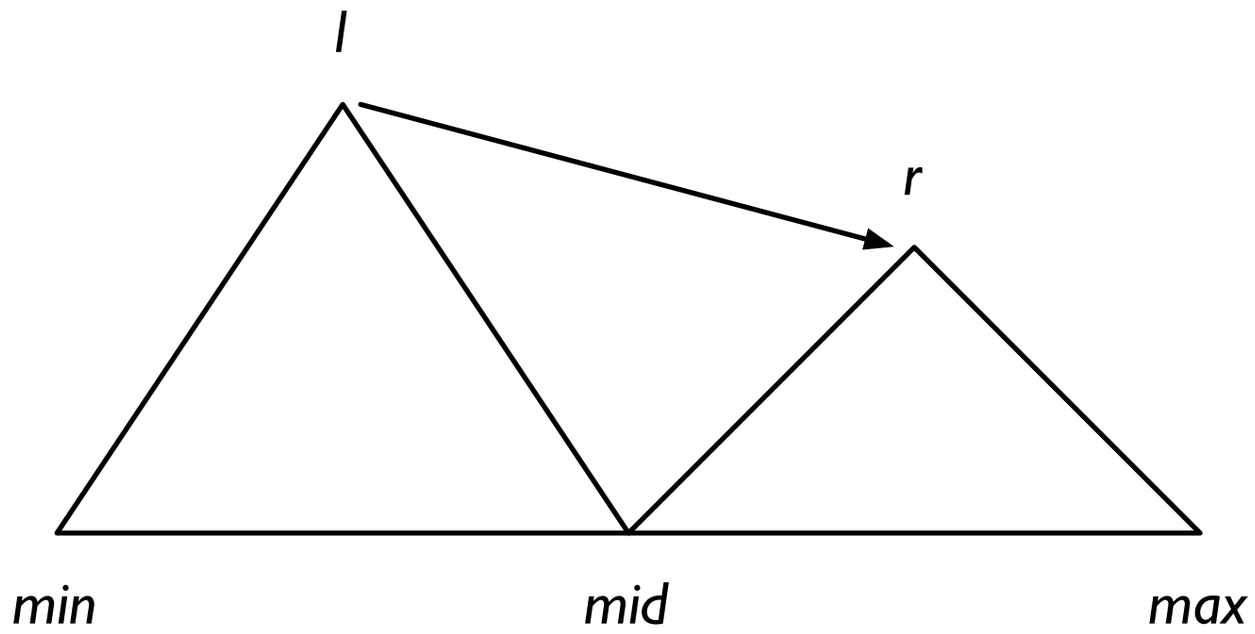
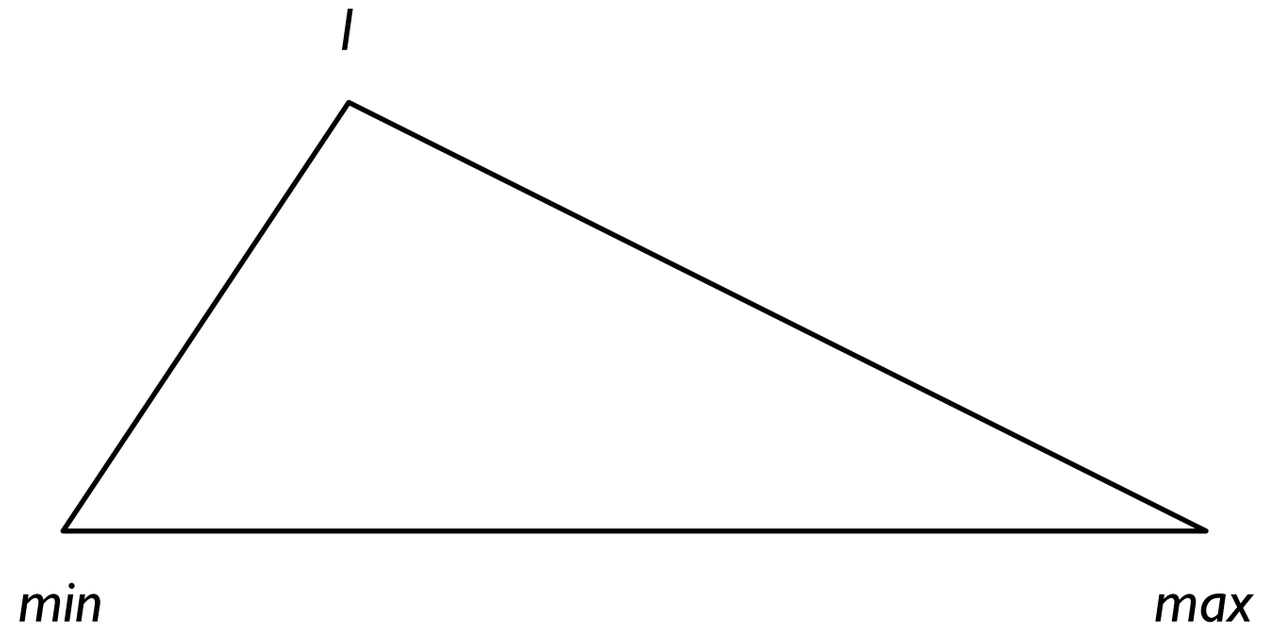
Basic idea



In Collins' algorithm, adding a left-to-right arc is done in one single step, specified by 5 positions.

Basic idea



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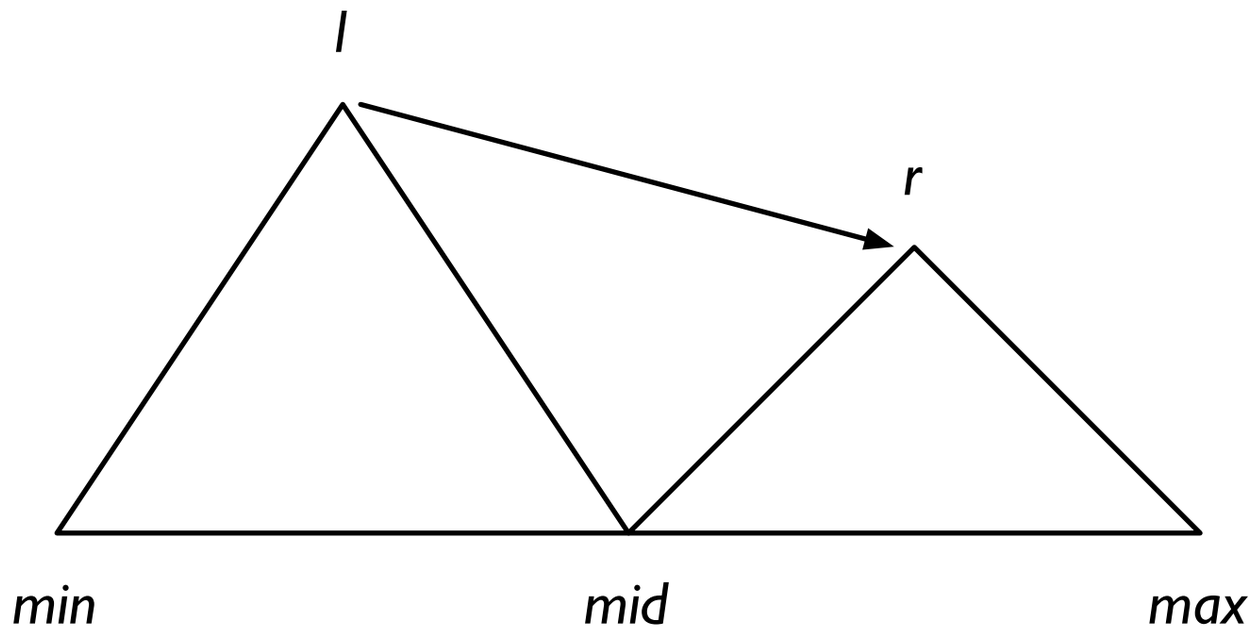
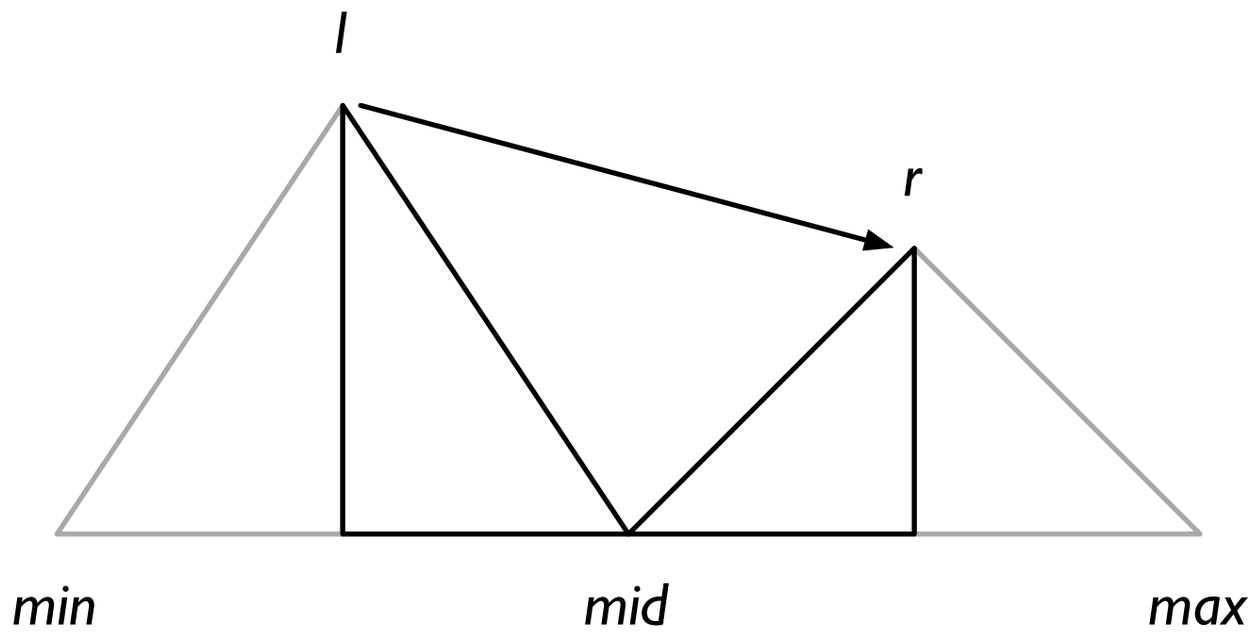


In Eisner's algorithm, the same thing is done in three steps, each one specified by 3 positions.

Basic idea



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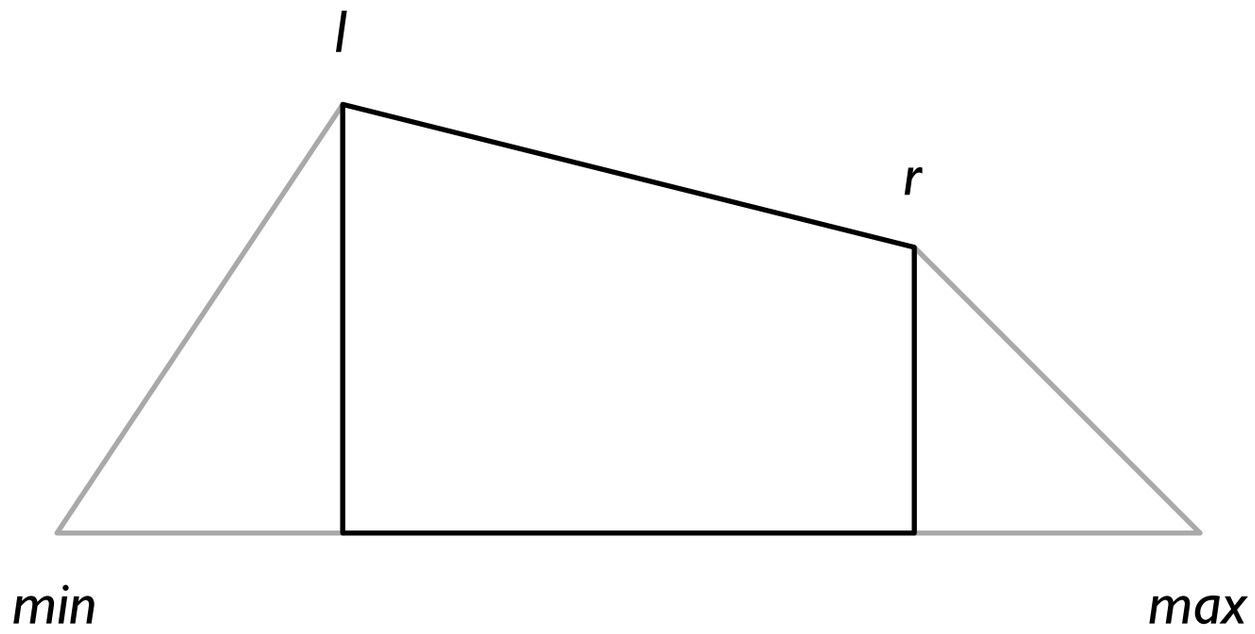
Eisner's algorithm

Basic idea



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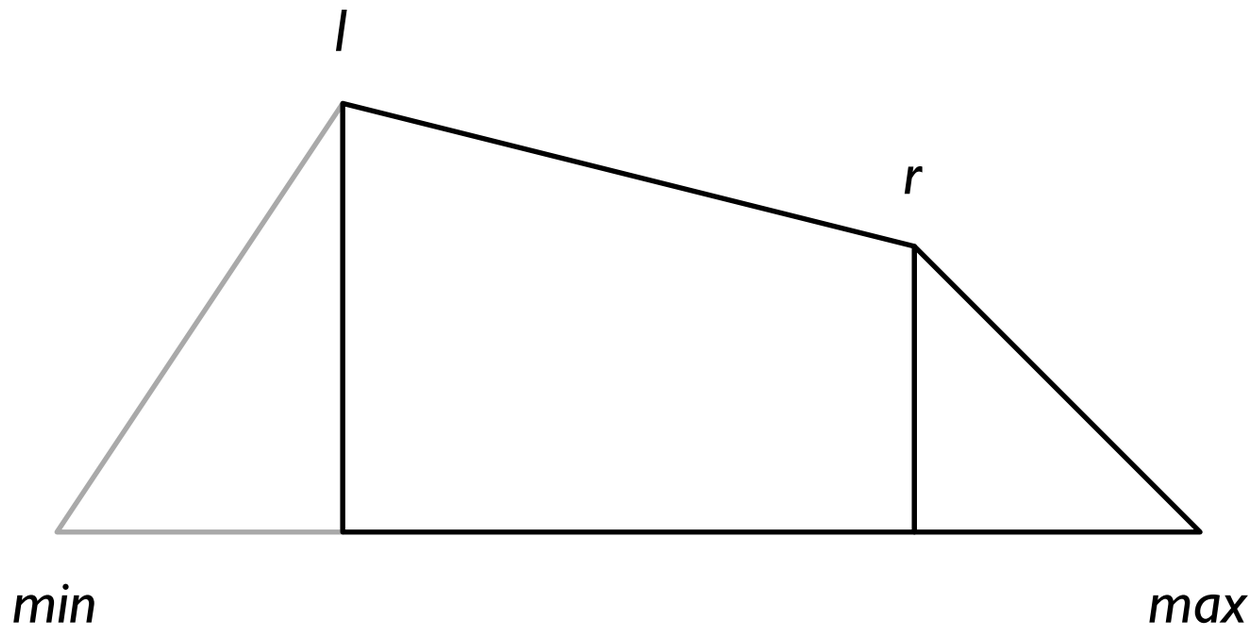

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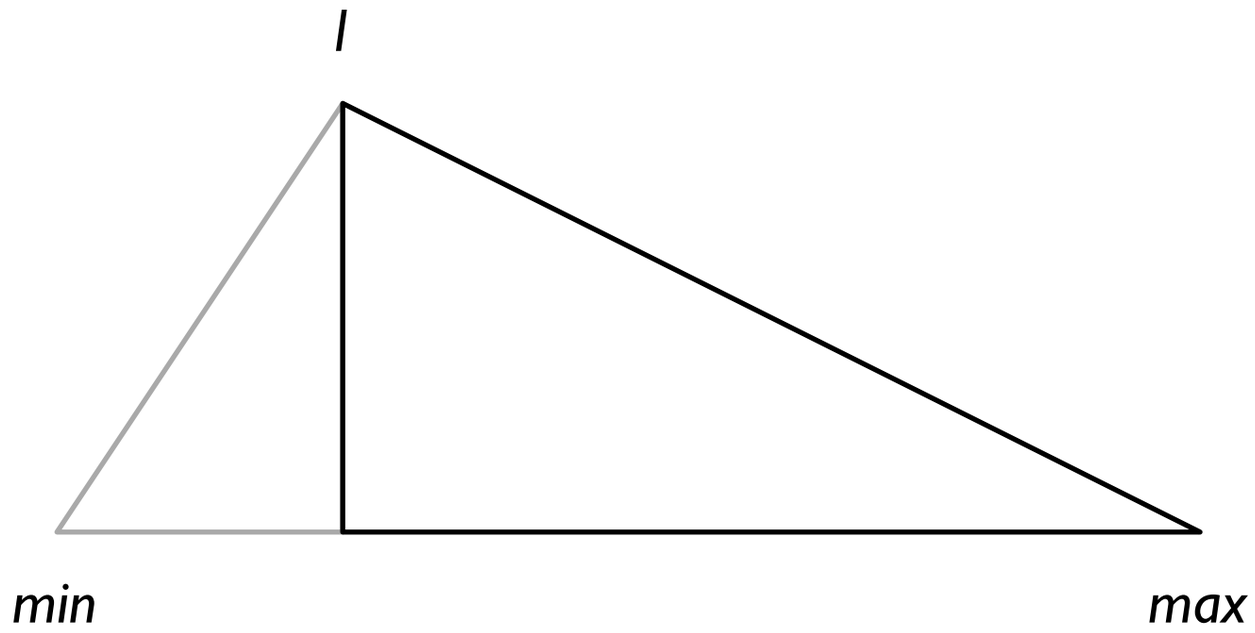

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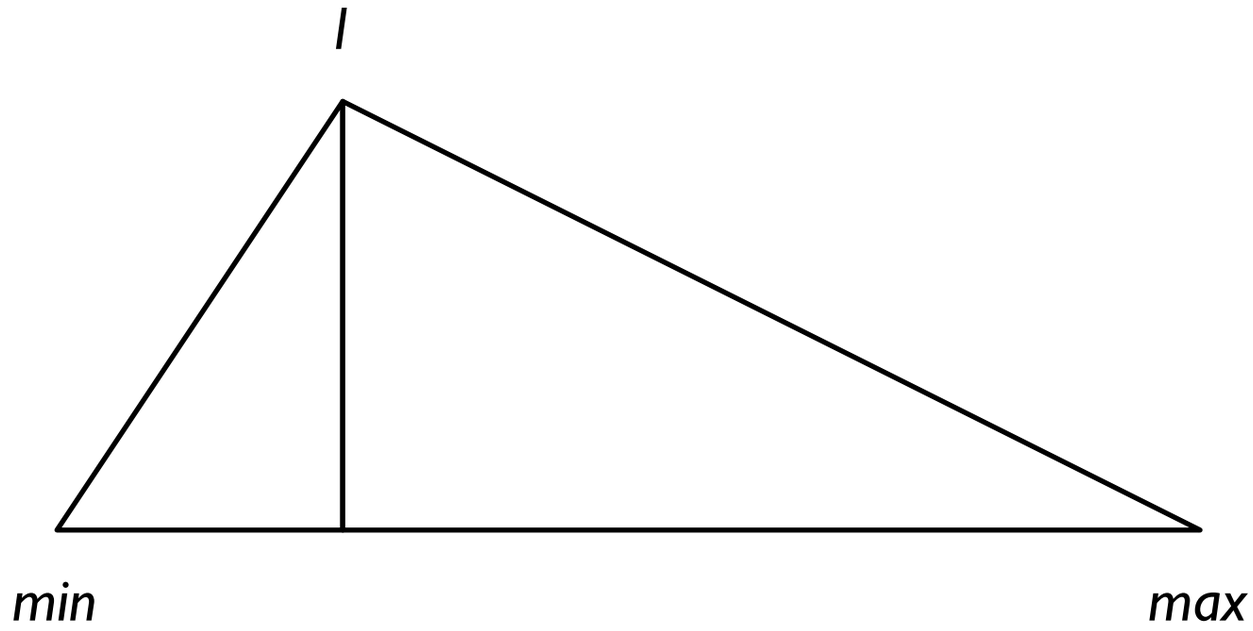

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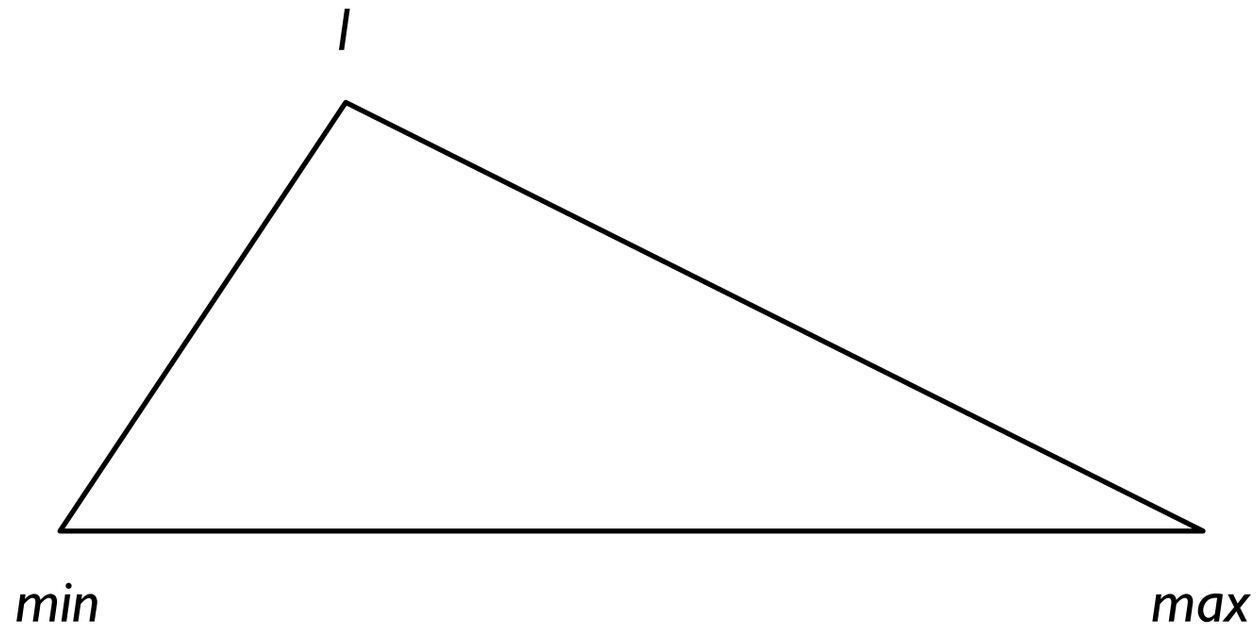
Eisner's algorithm

Basic idea



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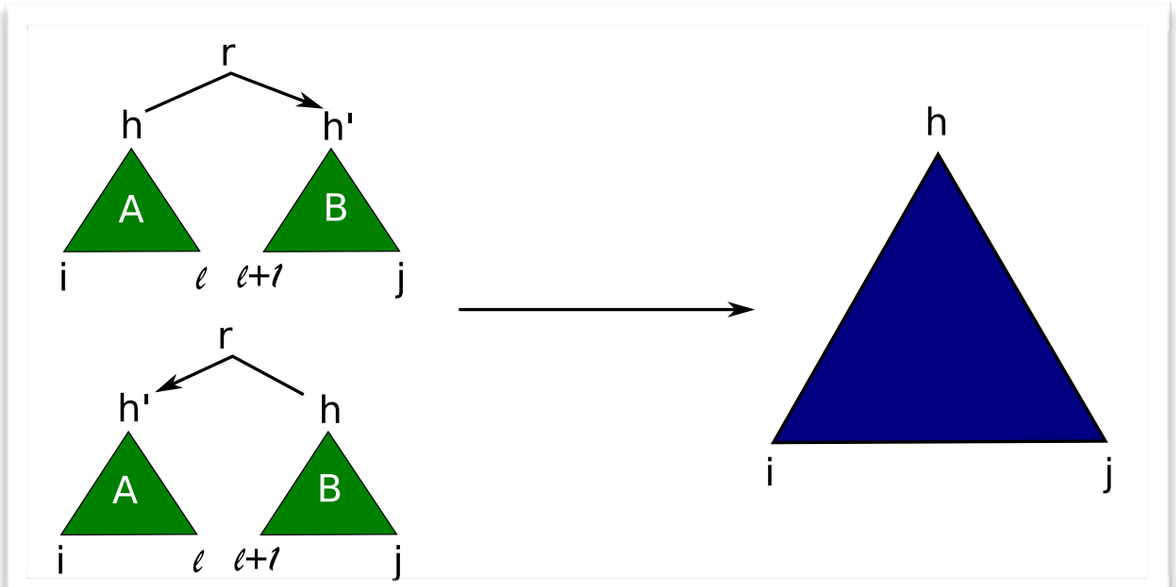
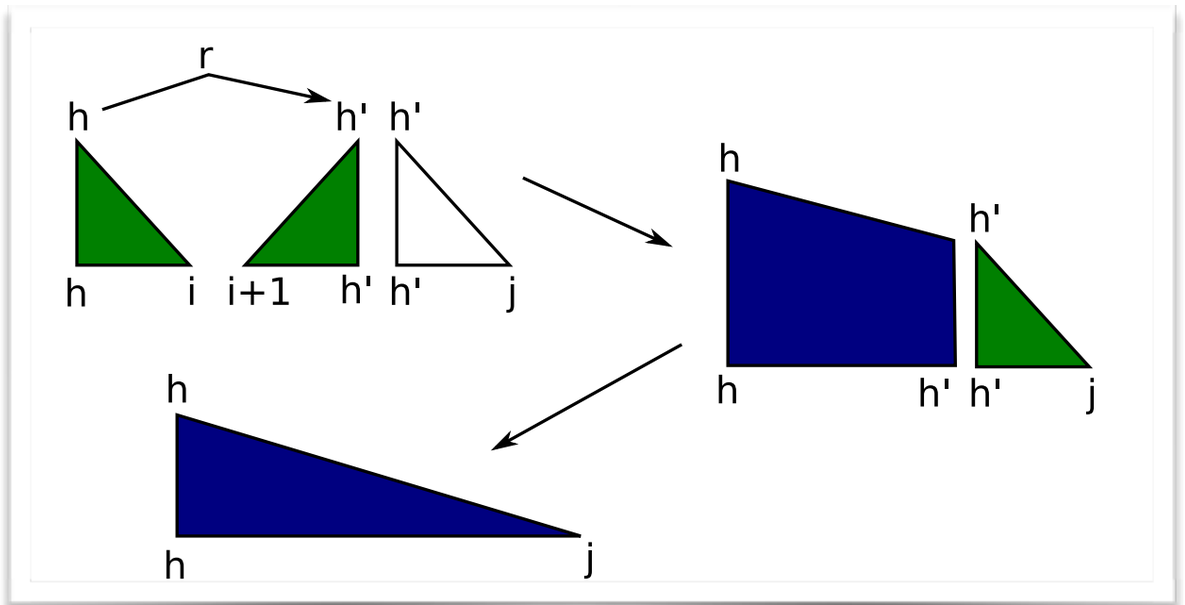
Comparison





Adding a left-to-right arc

```
for each [min, max] with max - min > 1 do
    double best- = score[min][max][left][-]
    double best+ = score[min][max][left][+]
    for each mid from min + 1 to max - 1 do
        t1 = score[min][mid][left][+]
        t2 = score[mid][max][right][+]
        double current- = t1 + t2 + score(min → max)
        if current- > best- then
            best- = current-
        t3 = score[min][mid+1][left][-]
        t4 = score[mid][max][left][+]
        double current+ = t3 + t4
        if current+ > best+ then
            best+ = current+
    score[min][max][left][-] = best-
    score[min][max][left][+] = best+
```

Adding a right-to-left arc

```
for each [min, max] with max - min > 1 do
    double best- = score[min][max][right][-]
    double best+ = score[min][max][right][+]
    for each mid from min + 1 to max - 1 do
        t1 = score[min][mid][left][+]
        t2 = score[mid][max][right][+]
        double current- = t1 + t2 + score(max → min)
        if current- > best- then
            best- = current-
        t3 = score[min][mid][right][+]
        t4 = score[mid-1][max][right][-]
        double current+ = t3 + t4
        if current+ > best+ then
            best+ = current+
    score[min][max][right][-] = best-
    score[min][max][right][+] = best+
```

Finishing up

- Find mid (from 1 to n) that maximizes:

  score[0][mid][right][+] + score[mid-1][n][left][+]

- Left and right half-trees are built independently, combined only in the last step

Summary

- Eisner's algorithm is an improvement over Collin's algorithm that runs in time $O(|w|^3)$.
- The same scoring model can be used.
- The same technique for extending the parser to labeled parsing can be used.
- Eisner's algorithm is the basis of current arc-factored dependency parsers.