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# Earley's algorithm

## Discussion

Syntactic analysis/parsing

2024

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Based on slides by Marco Kuhlmann



# Today

- Presentation of Earley's algorithm
- Working through an exercise on it
- If time:
  - Advanced PCFG



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# The Earley algorithm



# CKY versus Earley

- The CKY algorithm has two disadvantages:
  - It can only handle restricted grammars.
  - It does not use top–down information.
- The Earley algorithm does not have these:
  - It can handle arbitrary grammars.
  - It does use top–down information.
  - On the downside, it is more complicated.



# The algorithm

- Start with the start symbol  $S$ .
- Take the leftmost nonterminal and **predict** all possible expansions.
- If the next symbol in the expansion is a word, match it against the input sentence (**scan**); otherwise, repeat.
- If there is nothing more to expand, the subtree is **complete**; in this case, continue with the next incomplete subtree.



# Dotted rules

- A **dotted rule** is a partially processed rule.

*Example:*  $S \rightarrow NP \cdot VP$

- The dot can be placed in front of the first symbol, behind the last symbol, or between two symbols on the right-hand side of a rule.
- The general form of a dotted rule thus is  $A \rightarrow \alpha \cdot \beta$ , where  $A \rightarrow \alpha\beta$  is the original, non-dotted rule.



# Inference rules

<b>Axiom</b>	$[0, 0, S \rightarrow \cdot \alpha]$	$S \rightarrow \alpha$
<b>Predict</b>	$\frac{[i, j, A \rightarrow \alpha \cdot B \beta]}{[j, j, B \rightarrow \cdot \gamma]}$	$B \rightarrow \gamma$
<b>Scan</b>	$\frac{[i, j, A \rightarrow \alpha \cdot a \beta]}{[i, j+1, A \rightarrow \alpha a \cdot \beta]}$	$w_j = a$
<b>Complete</b>	$\frac{[i, j, A \rightarrow \alpha \cdot B \beta] \quad [j, k, B \rightarrow \gamma \cdot]}{[i, k, A \rightarrow \alpha B \cdot \beta]}$	



# Pseudo code I

```
function EARLEY-PARSE(words, grammar) returns chart  
  
ENQUEUE( $(\gamma \rightarrow \bullet S, [0, 0])$ , chart[0])  
for  $i \leftarrow$  from 0 to LENGTH(words) do  
  for each state in chart[ $i$ ] do  
    if INCOMPLETE?(state) and  
      NEXT-CAT(state) is not a part of speech then  
        PREDICTOR(state)  
    elseif INCOMPLETE?(state) and  
      NEXT-CAT(state) is a part of speech then  
        SCANNER(state)  
    else  
      COMPLETER(state)  
  end  
end  
return(chart)
```





## Pseudo code 2

```
procedure PREDICTOR( $(A \rightarrow \alpha \bullet B \beta, [i, j])$ )  
  for each  $(B \rightarrow \gamma)$  in GRAMMAR-RULES-FOR( $B, grammar$ ) do  
    ENQUEUE( $(B \rightarrow \bullet \gamma, [j, j])$ ,  $chart[j]$ )  
end  
  
procedure SCANNER( $(A \rightarrow \alpha \bullet B \beta, [i, j])$ )  
  if  $B \subset PARTS-OF-SPEECH(word[j])$  then  
    ENQUEUE( $(B \rightarrow word[j], [j, j+1])$ ,  $chart[j+1]$ )  
  
procedure COMPLETER( $(B \rightarrow \gamma \bullet, [j, k])$ )  
  for each  $(A \rightarrow \alpha \bullet B \beta, [i, j])$  in  $chart[j]$  do  
    ENQUEUE( $(A \rightarrow \alpha B \bullet \beta, [i, k])$ ,  $chart[k]$ )  
end
```



# Recogniser/parser

- When parsing is complete, is there a chart entry?  
[0, n,  $S \rightarrow \alpha \bullet$  ]
- Recognizer
- If we want a parser, we have to add back pointers, and retrieve a tree
- Earley's algorithm can be used for PCFGs, but it is more complicated than for CKY



# Coming up

- Monday, Feb. 26: Supervision
  - Assignment 3
- Wednesday, Feb 28: Lecture graph-based parsing
- Monday, March 4: Seminar 2



# Coming deadlines

- Project:
  - Choose individual or pair project
    - Sign up in Studium
  - Decide on your topic: information on the web page!
  - Proposal: February 26
- Project seminar moved to March 25, 9-12!
- Assignment 2: Feb. 22
- Assignment 3: March 11 (moved from March 4)



# Earley's algorithm as a project

- Suitable scope: implement a recognizer
  - Use a small toy grammar to show that it works
- A parser is a bit too much to code, but you are expected to discuss it in your report
- Also: read at least one relevant article (the math can be a bit challenging, so OK to read it at a higher level of abstraction)



# Missing a deadline

- If you miss a deadline (and you have not contacted Sara beforehand due to extraordinary circumstances)
- You can submit at a second deadline:
  - All assignments: April 1
  - Project and seminar report: April 15